| 分類／Classification | $\begin{aligned} & \square \mathrm{tDS} \quad \square \mathrm{tGW} \\ & \nabla \mathrm{I} / \mathrm{O} \text { Card } \end{aligned}$ | PETL／tET／tPET <br> VXC Card | DS／PDS／PPDSVxComm |  | tM－752N <br> Other |
| :---: | :---: | :---: | :---: | :---: | :---: |
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## Q：How to process data transmission from multiple DI／DO channels？

A：General Digital I／O cards have 8 channels per port，and each uses hexadecimal format，regardless of whether they are input or output．The following will discuss how the channel number corresponds to a hexadecimal value．

When an 8－Bit DO port is configured，the data format can be expressed using the following equation：

D／O byte $=00000000$（bit）$=0 \times 00$（hex）（see Appendix 1）．

It should be noted that the data information is read from left to right，respectively，i．e．，Bit 7 to Bit 0 ，where Bit 7 ～Bit 4 is combined as a group，and Bit 3 ～Bit 0 is combined as another group，as shown in the table below．Refer to the settings for the channels on each card for the channel definition of each port number．

|  | Bit 7～Bit 4 | Bit 3～Bit 0 |
| :---: | :---: | :---: |
| D／O Byte | 0000 | 0000 |

## －Operation Examples

## ＞Digital Output

Step 1：Output＂Bit 2＂：D／O＝ 00000100 （bit）＝0x04（hex）．
Step 2：Output＂Bit 5 and Bit 6 ＂：D／O $=01100100$（bit）$=0 \times 64$（hex）．
Step 3：Output＂Bit 7＂and＂Bit 2＂are not output：D／O＝ 11100000 （bit）＝0xEO（hex）．

## ＞Digital Input

1．Assuming the $\mathrm{D} / \mathrm{l}$ port receives the data set $0 \times 3 \mathrm{D}$ ．
＝＞D／I＝0x3D（hex）＝ 00111101 （bit）．
＝＞Indicates that Bits $0,2,3,4$ ，and 5 have a value，which is means that Input channels $0,2,3,4$ ，and 5 have data．
2. Assuming $D / I$ port receives the data set $0 \times 80$.
=> D/I = 0x80 (hex) = 10000000 (bit).
=> Indicates that only Bit 7 has a value, which means that only Input channel 7 has data.

Pros: When reading or writing data, multi-channel data transmission allows at least 8 channels of data to be processed simultaneously, which can effectively improve processing efficiency.

Cons: If only a value from a single channel is to be processed, value must be converted. Refer to the following examples for more details.

## - Binary Operation (for $\mathrm{C} / \mathrm{C}++$ )

Mask Computing:
Mask $=00000001$ left-shift 3 bits => 00001000
Mask = Invert Mask => 11110111
Mask Off for Bit N:

$$
\text { Mask }=\sim(1 \ll N)
$$

Result = Data \& Mask


Mask Computing:
Mask $=00000001$ left-shift 3 bits => 00001000
Set On for Bit N :
Mask = $1 \ll N$
Result = Data | Mask


Get Status of Bit 3 (Data $=1010$ 1010)
Data right-shift 3 bits $=00010101$
Data Mask Off with 00000001
00010101 \& $00000001=00000001$

Get Status of Bit N:

$$
\text { Result = (Data >> N) \& } 1
$$

## - Binary Operation (for VB)

Mask Computing:
Mask = 00000001 left-shift 3 bits => 00001000
Mask = Invert Mask => 11110111
Mask Off for Bit N :
Mask $=\operatorname{Not}\left(2^{\wedge} \mathrm{N}\right)$
Result = Data and Mask

Mask Computing:
Mask = 00000001 left-shift 3 bits => 00001000
Set On for Bit N:
Mask $=2$ ^ N
Result = Data or Mask

Get Status of Bit $3($ Data $=1010$ 1010)
Data right-shift 3 bits $=00010101$
Data Mask Off with 00000001
$00010101 \& 00000001=00000001$

Get Status of Bit N:
Result $=\left(\right.$ Data $\left.\backslash\left(2^{\wedge} \mathrm{N}\right)\right)$ and 1

## Set On for Bit 3

| Data | 1010 | 0010 |
| ---: | ---: | ---: |
| OR | 0000 | 1000 |
| Result | 1010 | 1010 |

Mask Off for Bit 3

| Data | 1010 | 1010 |
| ---: | ---: | ---: |
| AND | 1111 | 0111 |
| Result | 1010 | 0010 |

Appendix 1: Binary Transfer Hexadecimal

| Binary | Hexadecimal | Binary | Hexadecimal |
| :---: | :---: | :---: | :---: |
| 0000 | $0 \times 0$ | 1000 | $0 \times 8$ |
| 0001 | $0 \times 1$ | 1001 | $0 \times 9$ |
| 0010 | $0 \times 2$ | 1010 | $0 \times \mathrm{A}(10)$ |
| 0011 | $0 \times 3$ | 1011 | $0 \times \mathrm{B}(11)$ |
| 0100 | $0 \times 4$ | 1100 | $0 \times \mathrm{C}(12)$ |
| 0101 | $0 \times 5$ | 1101 | $0 \times \mathrm{D}(13)$ |
| 0110 | $0 \times 6$ | 1110 | $0 \times \mathrm{E}(14)$ |
| 0111 | $0 \times 7$ | 1111 | $0 \times \mathrm{F}(15)$ |

Appendix 2: Binary Operation (Bitwise)

| AND | 0 | 1 |
| :---: | :---: | :---: |
| 0 | 0 | 0 |
| 1 | 0 | 1 |


| OR | 0 | 1 |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 1 | 1 | 1 |



| NOT |  |
| :---: | :---: |
| 0 | 1 |
| 1 | 0 |


| XOR | 0 | 1 |
| :---: | :---: | :---: |
| 0 | 0 | 1 |
| 1 | 1 | 0 |



